Chapter 4: SQL

- Basic Structure
- Set Operations
- Aggregate Functions
- Null Values
- Nested Subqueries
- Derived Relations
- Views
- Modification of the Database
- Joined Relations
- Data Definition Language
- Embedded SQL, ODBC and JDBC
SQL is based on set and relational operations with certain modifications and enhancements.

A typical SQL query has the form:

```
select A_1, A_2, ..., A_n
from r_1, r_2, ..., r_m
where P
```

- \( A_i \)'s represent attributes
- \( r_i \)'s represent relations
- \( P \) is a predicate.

This query is equivalent to the relational algebra expression:

\[
\Pi_{A_1, A_2, ..., A_n}(\sigma_P (r_1 \times r_2 \times \ldots \times r_m))
\]

The result of an SQL query is a relation.
The select Clause

- The `select` clause list the attributes desired in the result of a query
  - corresponds to the projection operation of the relational algebra
- E.g. find the names of all branches in the `loan` relation
  ```sql
  select branch-name 
  from loan
  ```
- In the “pure” relational algebra syntax, the query would be:
  ```sql
  \Pi_{branch-name}(loan)
  ```
- **NOTE:** SQL does not permit the ‘-’ character in names,
  - Use, e.g., `branch_name` instead of `branch-name` in a real implementation.
  - We use ‘-’ since it looks nicer!
- **NOTE:** SQL names are case insensitive, i.e. you can use capital or small letters.
  - You may wish to use upper case where-ever we use bold font.
The select Clause (Cont.)

- SQL allows duplicates in relations as well as in query results.
- To force the elimination of duplicates, insert the keyword **distinct** after **select**.
- Find the names of all branches in the *loan* relations, and remove duplicates

  ```sql
  select distinct branch-name
  from loan
  ```

- The keyword **all** specifies that duplicates not be removed.

  ```sql
  select all branch-name
  from loan
  ```
The select Clause (Cont.)

- An asterisk in the select clause denotes “all attributes”
  
  ```
  select *
  from loan
  ```

- The `select` clause can contain arithmetic expressions involving the operation, `+`, `−`, `∗`, and `/`, and operating on constants or attributes of tuples.

- The query:
  
  ```
  select loan-number, branch-name, amount * 100
  from loan
  ```

  would return a relation which is the same as the `loan` relations, except that the attribute `amount` is multiplied by 100.
The where Clause

- The **where** clause specifies conditions that the result must satisfy.
  - It corresponds to the selection predicate of the relational algebra.

- To find all loan number for loans made at the Perryridge branch with loan amounts greater than $1200:

  ```sql
  select loan-number
  from loan
  where branch-name = 'Perryridge' and amount > 1200
  ```

- Comparison results can be combined using the logical connectives **and**, **or**, and **not**.

- Comparisons can be applied to results of arithmetic expressions.
SQL includes a **between** comparison operator.

E.g. Find the loan number of those loans with loan amounts between $90,000 and $100,000 (that is, \( \geq 90,000 \) and \( \leq 100,000 \))

```sql
select loan-number
from loan
where amount between 90000 and 100000
```
The from Clause

- The **from** clause lists the relations involved in the query
  - corresponds to the Cartesian product operation of the relational algebra.

- Find the Cartesian product \( \text{borrower} \times \text{loan} \)
  ```sql
  select *
  from borrower, loan
  ```

- Find the name, loan number and loan amount of all customers having a loan at the Perryridge branch.
  ```sql
  select customer-name, borrower.loan-number, amount
  from borrower, loan
  where borrower.loan-number = loan.loan-number and branch-name = 'Perryridge'
  ```
The Rename Operation

- The SQL allows renaming relations and attributes using the `as` clause:
  
  \[ \text{old-name as new-name} \]

- Find the name, loan number and loan amount of all customers; rename the column name `loan-number` as `loan-id`.

  \[
  \text{select customer-name, borrower.loan-number as loan-id, amount from borrower, loan where borrower.loan-number = loan.loan-number}
  \]
Tuple Variables

- Tuple variables are defined in the **from** clause via the use of the **as** clause.

- Find the customer names and their loan numbers for all customers having a loan at some branch.

  ```sql
  select customer-name, T.loan-number, S.amount
  from borrower as T, loan as S
  where T.loan-number = S.loan-number
  ```

- Find the names of all branches that have greater assets than some branch located in Brooklyn.

  ```sql
  select distinct T.branch-name
  from branch as T, branch as S
  where T.assets > S.assets and S.branch-city = 'Brooklyn'
  ```
String Operations

- SQL includes a string-matching operator for comparisons on character strings. Patterns are described using two special characters:
  - percent (%). The % character matches any substring.
  - underscore (_). The _ character matches any character.

- Find the names of all customers whose street includes the substring “Main”.
  
  ```sql
  select customer-name 
  from customer 
  where customer-street like ‘%Main%’ 
  ```

- Match the name “Main%”
  
  ```sql
  like ‘Main\%’ escape ‘\’
  ```

- SQL supports a variety of string operations such as
  - concatenation (using “||”)
  - converting from upper to lower case (and vice versa)
  - finding string length, extracting substrings, etc.
Ordering the Display of Tuples

- List in alphabetic order the names of all customers having a loan in Perryridge branch
  
  ```sql
  select distinct customer-name
  from borrower, loan
  where borrower loan-number = loan.loan-number and
  branch-name = 'Perryridge'
  order by customer-name
  ```

- We may specify `desc` for descending order or `asc` for ascending order, for each attribute; ascending order is the default.
  
  E.g. `order by customer-name desc`
In relations with duplicates, SQL can define how many copies of tuples appear in the result.

*Multiset* versions of some of the relational algebra operators – given multiset relations $r_1$ and $r_2$:

1. $\sigma_\theta(r_1)$: If there are $c_1$ copies of tuple $t_1$ in $r_1$, and $t_1$ satisfies selections $\sigma_\theta$, then there are $c_1$ copies of $t_1$ in $\sigma_\theta(r_1)$.

2. $\Pi_A(r_1)$: For each copy of tuple $t_1$ in $r_1$, there is a copy of tuple $\Pi_A(t_1)$ in $\Pi_A(r_1)$ where $\Pi_A(t_1)$ denotes the projection of the single tuple $t_1$.

3. $r_1 \times r_2$: If there are $c_1$ copies of tuple $t_1$ in $r_1$ and $c_2$ copies of tuple $t_2$ in $r_2$, there are $c_1 \times c_2$ copies of the tuple $t_1 \cdot t_2$ in $r_1 \times r_2$. 
Duplicates (Cont.)

Example: Suppose multiset relations \( r_1 (A, B) \) and \( r_2 (C) \) are as follows:

\[
\begin{align*}
  r_1 &= \{(1, a), (2, a)\} \\
  r_2 &= \{(2), (3), (3)\}
\end{align*}
\]

Then \( \Pi_B(r_1) \) would be \{\(a\), \(a\)\}, while \( \Pi_B(r_1) \times r_2 \) would be

\[
\{(a,2), (a,2), (a,3), (a,3), (a,3), (a,3)\}
\]

SQL duplicate semantics:

\[
\text{select } A_1, A_2, ..., A_n \\
\text{from } r_1, r_2, ..., r_m \\
\text{where } P
\]

is equivalent to the multiset version of the expression:

\[
\Pi_{A_1, A_2, ..., A_n}(\sigma_P (r_1 \times r_2 \times ... \times r_m))
\]
Set Operations

- The set operations \texttt{union}, \texttt{intersect}, and \texttt{except} operate on relations and correspond to the relational algebra operations $\cup$, $\cap$, $\setminus$.

- Each of the above operations automatically eliminates duplicates; to retain all duplicates use the corresponding multiset versions \texttt{union all}, \texttt{intersect all} and \texttt{except all}.

Suppose a tuple occurs $m$ times in $r$ and $n$ times in $s$, then, it occurs:

- $m + n$ times in $r$ \texttt{union all} $s$
- $\min(m,n)$ times in $r$ \texttt{intersect all} $s$
- $\max(0, m - n)$ times in $r$ \texttt{except all} $s$
Set Operations

- Find all customers who have a loan, an account, or both:
  
  \[
  \text{(select customer-name from depositor)}
  \text{union}
  \text{(select customer-name from borrower)}
  \]

- Find all customers who have both a loan and an account.
  
  \[
  \text{(select customer-name from depositor)}
  \text{intersect}
  \text{(select customer-name from borrower)}
  \]

- Find all customers who have an account but no loan.
  
  \[
  \text{(select customer-name from depositor)}
  \text{except}
  \text{(select customer-name from borrower)}
  \]
Aggregate Functions

- These functions operate on the multiset of values of a column of a relation, and return a value

  - **avg**: average value
  - **min**: minimum value
  - **max**: maximum value
  - **sum**: sum of values
  - **count**: number of values
Aggregate Functions (Cont.)

- Find the average account balance at the Perryridge branch.

  select avg (balance)
  from account
  where branch-name = 'Perryridge'

- Find the number of tuples in the customer relation.

  select count (*)
  from customer

- Find the number of depositors in the bank.

  select count (distinct customer-name)
  from depositor
Find the number of depositors for each branch.

```
select branch-name, count (distinct customer-name) 
from depositor, account 
where depositor.account-number = account.account-number 
group by branch-name
```

Note: Attributes in `select` clause outside of aggregate functions must appear in `group by` list.
Find the names of all branches where the average account balance is more than $1,200.

```sql
select branch-name, avg(balance)
from account
group by branch-name
having avg(balance) > 1200
```

Note: predicates in the **having** clause are applied after the formation of groups whereas predicates in the **where** clause are applied before forming groups.
Null Values

- It is possible for tuples to have a null value, denoted by `null`, for some of their attributes.
- `null` signifies an unknown value or that a value does not exist.
- The predicate `is null` can be used to check for null values.
  - E.g. Find all loan number which appear in the `loan` relation with null values for `amount`.
    
    
    ```sql
    select loan-number
    from loan
    where amount is null
    ```
  - The result of any arithmetic expression involving `null` is `null`.
    - E.g. 5 + null returns null.
  - However, aggregate functions simply ignore nulls.
    - more on this shortly.
Null Values and Three Valued Logic

- Any comparison with null returns unknown
  - E.g. 5 < null or null <> null or null = null

- Three-valued logic using the truth value unknown:
  - OR: (unknown or true) = true, (unknown or false) = unknown
    (unknown or unknown) = unknown
  - AND: (true and unknown) = unknown, (false and unknown) = false
    (unknown and unknown) = unknown
  - NOT: (not unknown) = unknown

- “P is unknown” evaluates to true if predicate P evaluates to unknown

- Result of where clause predicate is treated as false if it evaluates to unknown
Null Values and Aggregates

- Total all loan amounts

```sql
select sum (amount)
from loan
```

- Above statement ignores null amounts
- Result is null if there is no non-null amount
- All aggregate operations except `count(*)` ignore tuples with null values on the aggregated attributes.
Nested Subqueries

- SQL provides a mechanism for the nesting of subqueries.
- A subquery is a `select-from-where` expression that is nested within another query.
- A common use of subqueries is to perform tests for set membership, set comparisons, and set cardinality.
Example Query

- Find all customers who have both an account and a loan at the bank.

  ```sql
  select distinct customer-name
  from borrower
  where customer-name in (select customer-name
  from depositor)
  ```

- Find all customers who have a loan at the bank but do not have an account at the bank.

  ```sql
  select distinct customer-name
  from borrower
  where customer-name not in (select customer-name
  from depositor)
  ```
Example Query

Find all customers who have both an account and a loan at the Perryridge branch

```
select distinct customer-name
from borrower, loan
where borrower.loan-number = loan.loan-number and
branch-name = "Perryridge" and
(branch-name, customer-name) in
(select branch-name, customer-name
from depositor, account
where depositor.account-number =
account.account-number)
```

Note: Above query can be written in a much simpler manner. The formulation above is simply to illustrate SQL features.

(Schema used in this example)
Set Comparison

- Find all branches that have greater assets than some branch located in Brooklyn.

```sql
select distinct T.branch-name
from branch as T, branch as S
where T.assets > S.assets and
    S.branch-city = 'Brooklyn'
```

- Same query using `> some` clause

```sql
select branch-name
from branch
where assets > some
    (select assets
     from branch
     where branch-city = 'Brooklyn')
```
Definition of Some Clause

- \( F \text{ <comp> } \text{some } r \iff \exists t \in r \text{ s.t. } (F \text{ <comp> } t) \)

Where <comp> can be: \(<, \leq, >, =, \neq\)

\[
\begin{array}{c|c|c}
\text{0} & \text{5} & \text{false} \\
\hline
\text{5} & \text{6} & \text{true} \\
\end{array}
\]

(5 < some 5) = true
(read: 5 < some tuple in the relation)

(5 < some 5) = false

(5 = some 5) = true

(5 \neq some 5) = true (since 0 \neq 5)

(= some) \equiv \text{in}

However, (\neq some) \neq \text{not in}
Definition of all Clause

- \( F < \text{comp} > \text{all} \ r \iff \forall t \in r \ (F < \text{comp} > t) \)

  
  \[
  \begin{array}{c}
  0 \\
  5 \\
  6 \\
  \end{array}
  \]

  \((5 < \text{all} \begin{array}{c} \\
  5 \\
  6 \\
  \end{array}) = \text{false}\)

  
  \[
  \begin{array}{c}
  6 \\
  10 \\
  \end{array}
  \]

  \((5 < \text{all} \begin{array}{c} \\
  6 \\
  10 \\
  \end{array}) = \text{true}\)

  
  \[
  \begin{array}{c}
  4 \\
  5 \\
  \end{array}
  \]

  \((5 = \text{all} \begin{array}{c} \\
  4 \\
  5 \\
  \end{array}) = \text{false}\)

  
  \[
  \begin{array}{c}
  4 \\
  6 \\
  \end{array}
  \]

  \((5 \neq \text{all} \begin{array}{c} \\
  4 \\
  6 \\
  \end{array}) = \text{true} \) (since \(5 \neq 4 \) and \(5 \neq 6\))

\((\neq \text{all}) \equiv \text{not in}\)

However, \((= \text{all}) \neq \text{in}\)
Example Query

- Find the names of all branches that have greater assets than all branches located in Brooklyn.

```sql
select branch-name
from branch
where assets > all
 (select assets
from branch
where branch-city = 'Brooklyn')
```
Test for Empty Relations

- The **exists** construct returns the value **true** if the argument subquery is nonempty.
- **exists**  \( r \iff r \neq \emptyset \)
- **not exists**  \( r \iff r = \emptyset \)
Example Query

Find all customers who have an account at all branches located in Brooklyn.

```
select distinct S.customer-name
from depositor as S
where not exists ( 
    (select branch-name 
     from branch 
     where branch-city = 'Brooklyn')
except 
( select R.branch-name 
  from depositor as T, account as R 
  where T.account-number = R.account-number and 
    S.customer-name = T.customer-name))
```

(Schema used in this example)

- Note that $X - Y = \emptyset \iff X \subseteq Y$
- Note: Cannot write this query using $= \text{all}$ and its variants
The **unique** construct tests whether a subquery has any duplicate tuples in its result.

Find all customers who have at most one account at the Perryridge branch.

```sql
select T.customer-name
from depositor as T
where unique (
    select R.customer-name
    from account, depositor as R
    where T.customer-name = R.customer-name and
    R.account-number = account.account-number and
    account.branch-name = 'Perryridge')
```

**(Schema used in this example)**
Example Query

Find all customers who have at least two accounts at the Perryridge branch.

```
select distinct T.customer-name
from depositor T
where not unique (  
    select R.customer-name
    from account, depositor as R
    where T.customer-name = R.customer-name
    and
        R.account-number = account.account-number
    and
        account.branch-name = ‘Perryridge’)
```

(Schema used in this example)
Views

- Provide a mechanism to hide certain data from the view of certain users. To create a view we use the command:

```
create view v as <query expression>
```

where:

- `<query expression>` is any legal expression
- The view name is represented by `v`
Example Queries

- A view consisting of branches and their customers

  ```sql
  create view all-customer as
  (select branch-name, customer-name
   from depositor, account
   where depositor.account-number = account.account-number)
  union
  (select branch-name, customer-name
   from borrower, loan
   where borrower.loan-number = loan.loan-number)
  ```

- Find all customers of the Perryridge branch

  ```sql
  select customer-name
  from all-customer
  where branch-name = 'Perryridge'
  ```
Find the average account balance of those branches where the average account balance is greater than $1200.

```sql
select branch-name, avg-balance
from (select branch-name, avg(balance)
      from account
      group by branch-name)
      as result (branch-name, avg-balance)
      where avg-balance > 1200
```

Note that we do not need to use the `having` clause, since we compute the temporary (view) relation `result` in the `from` clause, and the attributes of `result` can be used directly in the `where` clause.
With clause allows views to be defined locally to a query, rather than globally. Analogous to procedures in a programming language.

Find all accounts with the maximum balance

```sql
with max-balance(value) as
  select max(balance)
  from account

select account-number
from account, max-balance
where account.balance = max-balance.value
```
Complex Query using With Clause

- Find all branches where the total account deposit is greater than the average of the total account deposits at all branches.

```
with branch-total (branch-name, value) as
    select branch-name, sum(balance) from account
    group by branch-name
with branch-total-avg(value) as
    select avg(value) from branch-total
select branch-name from branch-total, branch-total-avg
where branch-total.value >= branch-total-avg.value
```
Modification of the Database – Deletion

- Delete all account records at the Perryridge branch
  
  ```sql
  delete from account
  where branch-name = 'Perryridge'
  ```

- Delete all accounts at every branch located in Needham city.
  
  ```sql
  delete from account
  where branch-name in (select branch-name
                          from branch
                          where branch-city = 'Needham')
  
  delete from depositor
  where account-number in
    (select account-number
     from branch, account
     where branch-city = 'Needham'
     and branch.branch-name = account.branch-name)
  
  (Schema used in this example)
Delete the record of all accounts with balances below the average at the bank.

```
delete from account
    where balance < (select avg (balance)
                     from account)
```

Problem: as we delete tuples from deposit, the average balance changes

Solution used in SQL:
1. First, compute **avg** balance and find all tuples to delete
2. Next, delete all tuples found above (without recomputing **avg** or retesting the tuples)
Modification of the Database – Insertion

- Add a new tuple to `account`

  ```sql
  insert into account
  values ('A-9732', 'Perryridge', 1200)
  ```

  or equivalently

  ```sql
  insert into account (branch-name, balance, account-number)
  values ('Perryridge', 1200, 'A-9732')
  ```

- Add a new tuple to `account` with `balance` set to null

  ```sql
  insert into account
  values ('A-777', 'Perryridge', null)
  ```
Provide as a gift for all loan customers of the Perryridge branch, a $200 savings account. Let the loan number serve as the account number for the new savings account.

```sql
insert into account
select loan-number, branch-name, 200
from loan
where branch-name = 'Perryridge'

insert into depositor
select customer-name, loan-number
from loan, borrower
where branch-name = 'Perryridge'
    and loan.account-number = borrower.account-number
```

The select from where statement is fully evaluated before any of its results are inserted into the relation (otherwise queries like

```sql
insert into table1 select * from table1
```

would cause problems.
Increase all accounts with balances over $10,000 by 6%, all other accounts receive 5%.

Write two `update` statements:

```sql
update account
set balance = balance * 1.06
where balance > 10000

update account
set balance = balance * 1.05
where balance <= 10000
```

The order is important

Can be done better using the `case` statement (next slide)
Case Statement for Conditional Updates

- Same query as before: Increase all accounts with balances over $10,000 by 6%, all other accounts receive 5%.

```sql
update account
set balance = case
    when balance <= 10000 then balance * 1.05
    else balance * 1.06
end
```
Update of a View

- Create a view of all loan data in loan relation, hiding the amount attribute
  
  ```
  create view branch-loan as
  select branch-name, loan-number
  from loan
  ```

- Add a new tuple to branch-loan
  
  ```
  insert into branch-loan
  values ('Perryridge', 'L-307')
  ```

  This insertion must be represented by the insertion of the tuple
  
  ```
  ('L-307', 'Perryridge', null)
  ```

  into the loan relation

- Updates on more complex views are difficult or impossible to translate, and hence are disallowed.

- Most SQL implementations allow updates only on simple views (without aggregates) defined on a single relation
Transactions

- A transaction is a sequence of queries and update statements executed as a single unit
  - Transactions are started implicitly and terminated by one of
    - **commit work**: makes all updates of the transaction permanent in the database
    - **rollback work**: undoes all updates performed by the transaction.

- Motivating example
  - Transfer of money from one account to another involves two steps:
    - deduct from one account and credit to another
  - If one steps succeeds and the other fails, database is in an inconsistent state
  - Therefore, either both steps should succeed or neither should

- If any step of a transaction fails, all work done by the transaction can be undone by **rollback work**.

- Rollback of incomplete transactions is done automatically, in case of system failures
In most database systems, each SQL statement that executes successfully is automatically committed.

- Each transaction would then consist of only a single statement.
- Automatic commit can usually be turned off, allowing multi-statement transactions, but how to do so depends on the database system.
- Another option in SQL:1999: enclose statements within `begin atomic

...`

`end`
Joined Relations

- Join operations take two relations and return as a result another relation.
- These additional operations are typically used as subquery expressions in the `from` clause.
- Join condition – defines which tuples in the two relations match, and what attributes are present in the result of the join.
- Join type – defines how tuples in each relation that do not match any tuple in the other relation (based on the join condition) are treated.

<table>
<thead>
<tr>
<th>Join Types</th>
<th>Join Conditions</th>
</tr>
</thead>
<tbody>
<tr>
<td>inner join</td>
<td>natural</td>
</tr>
<tr>
<td>left outer join</td>
<td>on &lt;predicate&gt;</td>
</tr>
<tr>
<td>right outer join</td>
<td>using ((A_1, A_2, \ldots, A_n))</td>
</tr>
<tr>
<td>full outer join</td>
<td></td>
</tr>
</tbody>
</table>
**Joined Relations – Datasets for Examples**

- **Relation loan**

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
</tr>
</thead>
<tbody>
<tr>
<td>L-170</td>
<td>Downtown</td>
<td>3000</td>
</tr>
<tr>
<td>L-230</td>
<td>Redwood</td>
<td>4000</td>
</tr>
<tr>
<td>L-260</td>
<td>Perryridge</td>
<td>1700</td>
</tr>
</tbody>
</table>

- **Relation borrower**

<table>
<thead>
<tr>
<th>customer-name</th>
<th>loan-number</th>
</tr>
</thead>
<tbody>
<tr>
<td>Jones</td>
<td>L-170</td>
</tr>
<tr>
<td>Smith</td>
<td>L-230</td>
</tr>
<tr>
<td>Hayes</td>
<td>L-155</td>
</tr>
</tbody>
</table>

- **Note:** borrower information missing for L-260 and loan information missing for L-155
Joined Relations – Examples

- **loan inner join borrower on**
  \[ \text{loan.loan-number} = \text{borrower.loan-number} \]

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
<th>customer-name</th>
<th>loan-number</th>
</tr>
</thead>
<tbody>
<tr>
<td>L-170</td>
<td>Downtown</td>
<td>3000</td>
<td>Jones</td>
<td>L-170</td>
</tr>
<tr>
<td>L-230</td>
<td>Redwood</td>
<td>4000</td>
<td>Smith</td>
<td>L-230</td>
</tr>
</tbody>
</table>

- **loan left outer join borrower on**
  \[ \text{loan.loan-number} = \text{borrower.loan-number} \]

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
<th>customer-name</th>
<th>loan-number</th>
</tr>
</thead>
<tbody>
<tr>
<td>L-170</td>
<td>Downtown</td>
<td>3000</td>
<td>Jones</td>
<td>L-170</td>
</tr>
<tr>
<td>L-230</td>
<td>Redwood</td>
<td>4000</td>
<td>Smith</td>
<td>L-230</td>
</tr>
<tr>
<td>L-260</td>
<td>Perryridge</td>
<td>1700</td>
<td>null</td>
<td>null</td>
</tr>
</tbody>
</table>
### Joined Relations – Examples

- **loan natural inner join borrower**

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
<th>customer-name</th>
</tr>
</thead>
<tbody>
<tr>
<td>L-170</td>
<td>Downtown</td>
<td>3000</td>
<td>Jones</td>
</tr>
<tr>
<td>L-230</td>
<td>Redwood</td>
<td>4000</td>
<td>Smith</td>
</tr>
</tbody>
</table>

- **loan natural right outer join borrower**

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
<th>customer-name</th>
</tr>
</thead>
<tbody>
<tr>
<td>L-170</td>
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<td>3000</td>
<td>Jones</td>
</tr>
<tr>
<td>L-230</td>
<td>Redwood</td>
<td>4000</td>
<td>Smith</td>
</tr>
<tr>
<td>L-155</td>
<td>null</td>
<td>null</td>
<td>Hayes</td>
</tr>
</tbody>
</table>
Joined Relations – Examples

- **loan full outer join borrower using (loan-number)**

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
<th>customer-name</th>
</tr>
</thead>
<tbody>
<tr>
<td>L-170</td>
<td>Downtown</td>
<td>3000</td>
<td>Jones</td>
</tr>
<tr>
<td>L-230</td>
<td>Redwood</td>
<td>4000</td>
<td>Smith</td>
</tr>
<tr>
<td>L-260</td>
<td>Perryridge</td>
<td>1700</td>
<td>null</td>
</tr>
<tr>
<td>L-155</td>
<td>null</td>
<td>null</td>
<td>Hayes</td>
</tr>
</tbody>
</table>

- Find all customers who have either an account or a loan (but not both) at the bank.

  ```
  select customer-name
  from (depositor natural full outer join borrower)
  where account-number is null or loan-number is null
  ```
Data Definition Language (DDL)

Allows the specification of not only a set of relations but also information about each relation, including:

- The schema for each relation.
- The domain of values associated with each attribute.
- Integrity constraints
- The set of indices to be maintained for each relation.
- Security and authorization information for each relation.
- The physical storage structure of each relation on disk.
Domain Types in SQL

- **char(n)**. Fixed length character string, with user-specified length $n$.
- **varchar(n)**. Variable length character strings, with user-specified maximum length $n$.
- **int**. Integer (a finite subset of the integers that is machine-dependent).
- **smallint**. Small integer (a machine-dependent subset of the integer domain type).
- **numeric(p,d)**. Fixed point number, with user-specified precision of $p$ digits, with $n$ digits to the right of decimal point.
- **real, double precision**. Floating point and double-precision floating point numbers, with machine-dependent precision.
- **float(n)**. Floating point number, with user-specified precision of at least $n$ digits.

Null values are allowed in all the domain types. Declaring an attribute to be **not null** prohibits null values for that attribute.

**create domain** construct in SQL-92 creates user-defined domain types

```sql
create domain person-name char(20) not null
```
Date/Time Types in SQL (Cont.)

- **date.** Dates, containing a (4 digit) year, month and date
  - E.g. `date '2001-7-27'`
- **time.** Time of day, in hours, minutes and seconds.
  - E.g. `time '09:00:30'` `time '09:00:30.75'`
- **timestamp:** date plus time of day
  - E.g. `timestamp '2001-7-27 09:00:30.75'`
- **Interval:** period of time
  - E.g. `Interval '1' day`
  - Subtracting a date/time/timestamp value from another gives an interval value
  - Interval values can be added to date/time/timestamp values
- Can extract values of individual fields from date/time/timestamp
  - E.g. `extract (year from r.starttime)`
- Can cast string types to date/time/timestamp
  - E.g. `cast <string-valued-expression> as date`
Create Table Construct

- An SQL relation is defined using the `create table` command:

  ```sql
  create table r (A1 D1, A2 D2, ..., An Dn,
                  (integrity-constraint1),
                  ..., (integrity-constraintk))
  ``

  - \( r \) is the name of the relation
  - each \( A_i \) is an attribute name in the schema of relation \( r \)
  - \( D_i \) is the data type of values in the domain of attribute \( A_i \)

- Example:

  ```sql
  create table branch
  (branch-name char(15) not null,
   branch-city char(30),
   assets integer)
  ```
Integrity Constraints in Create Table

- **not null**
- **primary key** \((A_1, \ldots, A_n)\)
- **check** \((P)\), where \(P\) is a predicate

Example: Declare \(branch-name\) as the primary key for \(branch\) and ensure that the values of \(assets\) are non-negative.

```
create table branch
    (branch-name char(15),
     branch-city char(30),
     assets integer,
     primary key (branch-name),
     check (assets >= 0))
```

**primary key** declaration on an attribute automatically ensures **not null** in SQL-92 onwards, needs to be explicitly stated in SQL-89
The **drop table** command deletes all information about the dropped relation from the database.

The **alter table** command is used to add attributes to an existing relation.

\[
\text{alter table } r \text{ add } A \ D
\]

where \( A \) is the name of the attribute to be added to relation \( r \) and \( D \) is the domain of \( A \).

\( \checkmark \) All tuples in the relation are assigned \textit{null} as the value for the new attribute.

The **alter table** command can also be used to drop attributes of a relation

\[
\text{alter table } r \text{ drop } A
\]

where \( A \) is the name of an attribute of relation \( r \)

\( \checkmark \) Dropping of attributes not supported by many databases
The SQL standard defines embeddings of SQL in a variety of programming languages such as Pascal, PL/I, Fortran, C, and Cobol.

A language to which SQL queries are embedded is referred to as a host language, and the SQL structures permitted in the host language comprise embedded SQL.

The basic form of these languages follows that of the System R embedding of SQL into PL/I.

EXEC SQL statement is used to identify embedded SQL request to the preprocessor

EXEC SQL <embedded SQL statement > END-EXEC

Note: this varies by language. E.g. the Java embedding uses

```java
# SQL { .... } ;
```
Example Query

From within a host language, find the names and cities of customers with more than the variable amount dollars in some account.

- Specify the query in SQL and declare a cursor for it

EXEC SQL

```sql
declare c cursor for
select customer-name, customer-city
from depositor, customer, account
where depositor.customer-name = customer.customer-name
  and depositor.account-number = account.account-number
  and account.balance > :amount

END-EXEC
```
The **open** statement causes the query to be evaluated

```sql
EXEC SQL open c END-EXEC
```

The **fetch** statement causes the values of one tuple in the query result to be placed on host language variables.

```sql
EXEC SQL fetch c into :cn, :cc END-EXEC
```

Repeated calls to `fetch` get successive tuples in the query result.

- A variable called `SQLSTATE` in the SQL communication area (SQLCA) gets set to ‘02000’ to indicate no more data is available.

- The **close** statement causes the database system to delete the temporary relation that holds the result of the query.

```sql
EXEC SQL close c END-EXEC
```

Note: above details vary with language. E.g. the Java embedding defines Java iterators to step through result tuples.
Updates Through Cursors

- Can update tuples fetched by cursor by declaring that the cursor is for update

```sql
declare c cursor for
  select *
  from account
  where branch-name = 'Perryridge'
for update
```

- To update tuple at the current location of cursor

```sql
update account
set balance = balance + 100
where current of c
```
Dynamic SQL

- Allows programs to construct and submit SQL queries at runtime.
- Example of the use of dynamic SQL from within a C program.

```c
char * sqlprog = "update account
    set balance = balance * 1.05
    where account-number = ?"
EXEC SQL prepare dynprog from :sqlprog;
EXEC SQL execute dynprog using :account;
```

- The dynamic SQL program contains a ?, which is a place holder for a value that is provided when the SQL program is executed.
Open DataBase Connectivity (ODBC) standard

- Standard for application program to communicate with a database server.
- Application program interface (API) to
  - Open a connection with a database,
  - Send queries and updates,
  - Get back results.

Applications such as GUI, spreadsheets, etc. can use ODBC.
Each database system supporting ODBC provides a "driver" library that must be linked with the client program.

When client program makes an ODBC API call, the code in the library communicates with the server to carry out the requested action, and fetch results.

ODBC program first allocates an SQL environment, then a database connection handle.

Opens database connection using SQLConnect(). Parameters for SQLConnect:

- connection handle,
- the server to which to connect
- the user identifier,
- password

Must also specify types of arguments:

- SQL_NTS denotes previous argument is a null-terminated string.
int ODBCexample()
{
    RETCODE error;
    HENV env; /* environment */
    HDBC conn; /* database connection */
    SQLAllocEnv(&env);
    SQLAllocConnect(env, &conn);
    SQLConnect(conn, "aura.bell-labs.com", SQL_NTS, "avi", SQL_NTS, "avipassword", SQL_NTS);
    { .... Do actual work ... }
    SQLDisconnect(conn);
    SQLFreeConnect(conn);
    SQLFreeEnv(env);
}
Program sends SQL commands to the database by using SQLExecDirect

Result tuples are fetched using SQLFetch()

SQLBindCol() binds C language variables to attributes of the query result

- When a tuple is fetched, its attribute values are automatically stored in corresponding C variables.

- Arguments to SQLBindCol()
  - ODBC stmt variable, attribute position in query result
  - The type conversion from SQL to C.
  - The address of the variable.
  - For variable-length types like character arrays,
    - The maximum length of the variable
    - Location to store actual length when a tuple is fetched.
    - Note: A negative value returned for the length field indicates null value

Good programming requires checking results of every function call for errors; we have omitted most checks for brevity.
Main body of program

```
char branchname[80];
float balance;
int lenOut1, lenOut2;
HSTMT stmt;

SQLAllocStmt(conn, &stmt);
char * sqlquery = "select branch_name, sum(balance) from account group by branch_name";
error = SQLExecDirect(stmt, sqlquery, SQL_NTS);
if (error == SQL_SUCCESS) {
    SQLBindCol(stmt, 1, SQL_C_CHAR, branchname, 80, &lenOut1);
    SQLBindCol(stmt, 2, SQL_C_FLOAT, &balance, 0, &lenOut2);
    while (SQLFetch(stmt) >= SQL_SUCCESS) {
        printf(" %s  %g\n", branchname, balance);
    }
}
SQLFreeStmt(stmt, SQL_DROP);
```
More ODBC Features

- Prepared Statement
  - SQL statement prepared: compiled at the database
  - Can have placeholders: E.g. `insert into account values(?,?,?)`
  - Repeatedly executed with actual values for the placeholders

- Metadata features
  - Finding all the relations in the database and
  - Finding the names and types of columns of a query result or a relation in the database.

- By default, each SQL statement is treated as a separate transaction that is committed automatically.
  - Can turn off automatic commit on a connection
    - `SQLSetConnectOption(conn, SQL_AUTOCOMMIT, 0)`
  - Transactions must then be committed or rolled back explicitly by
    - `SQLTransact(conn, SQL_COMMIT)` or
    - `SQLTransact(conn, SQL_ROLLBACK)`
ODBC Conformance Levels

- Conformance levels specify subsets of the functionality defined by the standard.
  - Core
  - Level 1 requires support for metadata querying
  - Level 2 requires ability to send and retrieve arrays of parameter values and more detailed catalog information.

- SQL Call Level Interface (CLI) standard similar to ODBC interface, but with some minor differences.
JDBC

- JDBC is a Java API for communicating with database systems supporting SQL
- JDBC supports a variety of features for querying and updating data, and for retrieving query results
- JDBC also supports metadata retrieval, such as querying about relations present in the database and the names and types of relation attributes
- Model for communicating with the database:
  - Open a connection
  - Create a “statement” object
  - Execute queries using the Statement object to send queries and fetch results
  - Exception mechanism to handle errors
public static void JDBCexample(String dbid, String userid, String passwd) {
    try {
        Class.forName("oracle.jdbc.driver.OracleDriver");
        Connection conn = DriverManager.getConnection("jdbc:oracle:thin:@aura.bell-labs.com:2000:bankdb", userid, passwd);
        Statement stmt = conn.createStatement();
        ... Do Actual Work ....
        stmt.close();
        conn.close();
    }
    catch (SQLException sqle) {
        System.out.println("SQLException : " + sqle);
    }
}
Update to database

```java
try {
    stmt.executeUpdate("insert into account values ('A-9732', 'Perryridge', 1200)" watching); using SQL UPDATE statement.
}
catch (SQLException sqle) {
    System.out.println("Could not insert tuple. " + sqle);
}
```

Execute query and fetch and print results

```java
ResultSet rset = stmt.executeQuery("select branch_name, avg(balance) from account group by branch_name");
```
```java
while (rset.next()) {
    System.out.println(rset.getString("branch_name") + " " + rset.getFloat(2));
}
```
JDBC Code Details

- **Getting result fields:**
  
  rs.getString("branchname") and rs.getString(1) equivalent if branchname is the first argument of select result.

- **Dealing with Null values**
  
  ```java
  int a = rs.getInt("a");
  if (rs.wasNull()) System.out.println("Got null value");
  ```
Prepared Statement

- Prepared statement allows queries to be compiled and executed multiple times with different arguments

```java
PreparedStatement pStmt = conn.prepareStatement("insert into account values(?,?,?)");
pStmt.setString(1, "A-9732");
pStmt.setString(2, "Perryridge");
pStmt.setInt(3, 1200);
pStmt.executeUpdate();
pStmt.setString(1, "A-9733");
pStmt.executeUpdate();
```

- Beware: If value to be stored in database contains a single quote or other special character, prepared statements work fine, but creating a query string and executing it directly would result in a syntax error!
Other SQL Features

- SQL sessions
  - client *connects* to an SQL server, establishing a session
  - executes a series of statements
  - *disconnects* the session
  - can *commit* or *rollback* the work carried out in the session

- An SQL environment contains several components, including a user identifier, and a *schema*, which identifies which of several schemas a session is using.
Three-level hierarchy for naming relations.

- Database contains multiple catalogs
- Each catalog can contain multiple schemas
- SQL objects such as relations and views are contained within a schema

- e.g. catalog5.bank-schema.account

- Each user has a default catalog and schema, and the combination is unique to the user.
- Default catalog and schema are set up for a connection
- Catalog and schema can be omitted, defaults are assumed
- Multiple versions of an application (e.g. production and test) can run under separate schemas
Procedural Extensions and Stored Procedures

- SQL provides a **module** language
  - permits definition of procedures in SQL, with if-then-else statements, for and while loops, etc.
  - more in Chapter 9

- **Stored Procedures**
  - Can store procedures in the database
  - then execute them using the **call** statement
  - permit external applications to operate on the database without knowing about internal details

- These features are covered in Chapter 9 (Object Relational Databases)
Extra Material on JDBC and Application Architectures
Transactions in JDBC

- As with ODBC, each statement gets committed automatically in JDBC
- To turn off auto commit use
  ```java
  conn.setAutoCommit(false);
  ```
- To commit or abort transactions use
  ```java
  conn.commit()  or  conn.rollback()
  ```
- To turn auto commit on again, use
  ```java
  conn.setAutoCommit(true);
  ```
JDBC provides a class CallableStatement which allows SQL stored procedures/functions to be invoked.

```java
CallableStatement cs1 = conn.prepareCall( "{call proc (?,?)}" ) ;
CallableStatement cs2 = conn.prepareCall( "{? = call func (?,?)}" ) ;
```
The class ResultSetMetaData provides information about all the columns of the ResultSet.

Instance of this class is obtained by getMetaData( ) function of ResultSet.

Provides Functions for getting number of columns, column name, type, precision, scale, table from which the column is derived etc.

```java
ResultSetMetaData rsmd = rs.getMetaData( );
for ( int i = 1; i <= rsmd.getColumnCount( ); i++ ) {
    String name = rsmd.getColumnName(i);
    String typeName = rsmd.getColumnTypeName(i);
}
```
The class DatabaseMetaData provides information about database relations. It has functions for getting all tables, all columns of the table, primary keys, etc. For example, to print column names and types of a relation:

```java
DatabaseMetaData dbmd = conn.getMetaData( );
ResultSet rs = dbmd.getColumns( null, "BANK-DB", "account", "%" );
// Arguments: catalog, schema-pattern, table-pattern, column-pattern
// Returns: 1 row for each column, with several attributes such as
//          COLUMN_NAME, TYPE_NAME, etc.
while ( rs.next( ) ) {
    System.out.println( rs.getString("COLUMN_NAME") ,
        rs.getString("TYPE_NAME");
}
```

There are also functions for getting information such as:
- Foreign key references in the schema
- Database limits like maximum row size, maximum no. of connections, etc.
Applications can be built using one of two architectures

- Two tier model
  - Application program running at user site directly uses JDBC/ODBC to communicate with the database

- Three tier model
  - Users/programs running at user sites communicate with an application server. The application server in turn communicates with the database
Two-tier Model

- E.g. Java code runs at client site and uses JDBC to communicate with the backend server

- Benefits:
  - flexible, need not be restricted to predefined queries

- Problems:
  - Security: passwords available at client site, all database operation possible
  - More code shipped to client
  - Not appropriate across organizations, or in large ones like universities
Three Tier Model

CGI Program

Application/HTTP Server

Servlets

Database Server

JDBC

HTTP/Application Specific Protocol

Network

Client

Client

Client
Three-tier Model (Cont.)

- E.g. Web client + Java Servlet using JDBC to talk with database server
- Client sends request over http or application-specific protocol
- Application or Web server receives request
- Request handled by CGI program or servlets
- Security handled by application at server
  - Better security
  - Fine granularity security
- Simple client, but only packaged transactions
End of Chapter
### The loan and borrower Relations

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
<th>customer-name</th>
<th>loan-number</th>
</tr>
</thead>
<tbody>
<tr>
<td>L-170</td>
<td>Downtown</td>
<td>3000</td>
<td>Jones</td>
<td>L-170</td>
</tr>
<tr>
<td>L-230</td>
<td>Redwood</td>
<td>4000</td>
<td>Smith</td>
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</tr>
<tr>
<td>L-260</td>
<td>Perryridge</td>
<td>1700</td>
<td>Hayes</td>
<td>L-155</td>
</tr>
</tbody>
</table>
The Result of `loan` inner join `borrower` on `loan.loan-number = borrower.loan-number`
The Result of \textit{loan} left outer join \textit{borrower on loan-number}

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
<th>customer-name</th>
<th>loan-number</th>
</tr>
</thead>
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</tr>
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<td>L-260</td>
<td>Perryridge</td>
<td>1700</td>
<td>null</td>
<td>null</td>
</tr>
</tbody>
</table>
The Result of *loan* natural inner join *borrower*

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
<th>customer-name</th>
</tr>
</thead>
<tbody>
<tr>
<td>L-170</td>
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<tr>
<td>L-230</td>
<td>Redwood</td>
<td>4000</td>
<td>Smith</td>
</tr>
</tbody>
</table>
Join Types and Join Conditions

<table>
<thead>
<tr>
<th>Join types</th>
<th>Join Conditions</th>
</tr>
</thead>
<tbody>
<tr>
<td>inner join</td>
<td><strong>natural</strong></td>
</tr>
<tr>
<td>left outer join</td>
<td><strong>on</strong> &lt; predicate&gt;</td>
</tr>
<tr>
<td>right outer join</td>
<td><strong>using</strong> ( \bar{A}_1, A_1, \ldots, A_n )</td>
</tr>
<tr>
<td>full outer join</td>
<td></td>
</tr>
</tbody>
</table>
The Result of \textit{loan} natural right outer join \textit{borrower}

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
<th>customer-name</th>
</tr>
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<td>Redwood</td>
<td>4000</td>
<td>Smith</td>
</tr>
<tr>
<td>L-155</td>
<td>\textit{null}</td>
<td>\textit{null}</td>
<td>Hayes</td>
</tr>
</tbody>
</table>
The Result of \textit{loan} full outer join \textit{borrower} using(\textit{loan-number})

<table>
<thead>
<tr>
<th>loan-number</th>
<th>branch-name</th>
<th>amount</th>
<th>customer-name</th>
</tr>
</thead>
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</tr>
<tr>
<td>L-155</td>
<td>null</td>
<td>null</td>
<td>Hayes</td>
</tr>
</tbody>
</table>
create table customer
  (customer-name char(20),
   customer-street char(30),
   customer-city char(30),
   primary key (customer-name))

create table branch
  (branch-name char(15),
   branch-city char(30),
   assets integer,
   primary key (branch-name),
   check (assets >= 0))

create table account
  (account-number char(10),
   branch-name char(15),
   balance integer,
   primary key (account-number),
   check (balance >= 0))

create table depositor
  (customer-name char(20),
   account-number char(10),
   primary key (customer-name, account-number))